

section setdefs

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The ‘definition’ parts of this section are a formal specification of those parts of the mathematical tool-kit which do not depend on numbers. All objects are defined before they are used, and all are defined explicitly.

section *setdefs*

1. Sets

1.0.0.1. Name

\neq	–	Inequality
\notin	–	Non-membership

1.0.0.2. Definition

relation(\neq)

$$\neq [X] == \{x, y : X \mid \neg (x = y)\}$$

relation(\notin)

$$\notin [X] == \{x : X; S : \mathbb{P} X \mid \neg (x \in S)\}$$



Go Back

Full Screen

Close

Quit

1.0.0.3. Name

- \emptyset – Empty set
- \subseteq – Subset relation
- \subset – Proper subset relation
- \mathbb{P}_1 – Non-empty subsets

1.0.0.4. Definition

$$\emptyset[X] == \{x : X \mid false\}$$

$$relation(_ \subseteq _)$$

$$_ \subseteq _[X] == \{S, T : \mathbb{P} X \mid (\forall x : S \bullet x \in T)\}$$

$$relation(_ \subset _)$$

$$_ \subset _[X] == \{S, T : \mathbb{P} X \mid S \subseteq T \wedge S \neq T\}$$

$$\mathbb{P}_1 X == \{a : \mathbb{P} X \mid a \neq \emptyset\}$$

1.0.0.5. Name

- \cup – Set union
- \cap – Set intersection
- \setminus – Set difference
- \ominus – Set symmetric difference

1.0.0.6. Definition

function 30 leftassoc ($_ \cup _$)

$$_ \cup _[X] == \lambda S, T : \mathbb{P} X \bullet \{x : X \mid x \in S \vee x \in T\}$$

function 40 leftassoc ($_ \cap _$)

$$_ \cap _[X] == \lambda S, T : \mathbb{P} X \bullet \{x : X \mid x \in S \wedge x \in T\}$$

function 30 leftassoc ($_ \setminus _$)

$$_ \setminus _[X] == \lambda S, T : \mathbb{P} X \bullet \{x : X \mid x \in S \wedge x \notin T\}$$

function 25 leftassoc ($_ \ominus _$)

$$_ \ominus _[X] == \lambda S, T : \mathbb{P} X \bullet \{x : X \mid \neg (x \in S \Leftrightarrow x \in T)\}$$

1.0.0.7. Description These four functions are the only possible non-degenerate functions between two sets using only the information that they are of the same type as each other and that the result is of that same common type. The four functions form a closed family, in that if we take any two starting sets of the same type, and form any expression using any number of occurrences of those two sets and these four functions, the result will be one which could be expressed by a single use of one of these functions.

1.0.0.8. Name \bigcup – Generalized union
 \bigcap – Generalized intersection

$$\bigcup[X] == \lambda A : \mathbb{P}(\mathbb{P} X) \bullet \{x : X \mid (\exists S : A \bullet x \in S)\}$$

$$\bigcap[X] == \lambda A : \mathbb{P}(\mathbb{P} X) \bullet \{x : X \mid (\forall S : A \bullet x \in S)\}$$

1.0.0.9. Name \mathbb{F} – Finite sets
 \mathbb{F}_1 – Non-empty finite sets

$$\text{generic}(\mathbb{F} _)$$

$$\mathbb{F} X == \bigcap \{A : \mathbb{P}(\mathbb{P} X) \mid \emptyset \in A \wedge \forall a : A; x : X \bullet a \cup \{x\} \in A\}$$

$$\mathbb{F}_1 X == \mathbb{F} X \setminus \{\emptyset\}$$

$$\text{relation}(\text{finite } _)$$

$$\text{finite } _[X] == \mathbb{F} X$$

1.0.0.10. Description A subset S of X is finite if and only if it is a member of the set of finite subsets of X , $S \in \mathbb{F} X$. The finite subsets of X form the smallest set which contains the empty set and is closed under the action of adding single elements of X . The non-empty finite subsets of X form the smallest set which contains the singleton sets of X and is closed under the action of adding single elements of X . The sets in $\mathbb{F}_1 X$ are the non-empty members of $\mathbb{F} X$.

1.0.0.11. Name id – Identity relation

1.0.0.12. Definition

$\text{generic}(id _)$

$id\ X == \lambda x : X \bullet x$

2. Relations

2.0.0.13. Name

\leftrightarrow	– Binary relations
\mapsto	– Maplet
first, second	– Projection functions for ordered pairs

2.0.0.14. Definition

$\text{generic } 5 \text{ rightassoc } (_ \leftrightarrow _)$

$X \leftrightarrow Y == \mathbb{P}(X \times Y)$

$\text{function } 10 \text{ leftassoc } (_ \mapsto _)$

$_ \mapsto _[X, Y] == \lambda p : X \times Y \bullet p$

$$first[X, Y] == \lambda p : X \times Y \bullet p.1$$

$$second[X, Y] == \lambda p : X \times Y \bullet p.2$$

2.0.0.15. Name dom, ran – Domain and range of a relation

2.0.0.16. Definition

$$dom[X, Y] == \lambda R : X \leftrightarrow Y \bullet \{p : R \bullet p.1\}$$

$$ran[X, Y] == \lambda R : X \leftrightarrow Y \bullet \{p : R \bullet p.2\}$$

2.0.0.17. Name \triangleleft – Domain restriction
 \triangleright – Range restriction

2.0.0.18. Definition

function 61 rightassoc ($_ \triangleleft _$)

$$_ \triangleleft _[X, Y] == \lambda S : \mathbb{P} X; R : X \leftrightarrow Y \bullet \{p : R \mid p.1 \in S\}$$

function 60 leftassoc ($_ \triangleright _$)

$$_ \triangleright _[X, Y] == \lambda R : X \leftrightarrow Y; T : \mathbb{P} Y \bullet \{p : R \mid p.2 \in T\}$$

2.0.0.19. Name \triangleleft – Domain subtraction
 \triangleright – Range subtraction

2.0.0.20. Definition

function 61 rightassoc ($_ \triangleleft _$)

$$_ \triangleleft _[X, Y] == \lambda S : \mathbb{P} X; R : X \leftrightarrow Y \bullet \{p : R \mid p.1 \notin S\}$$

function 60 leftassoc ($_ \triangleright _$)

$$_ \triangleright _[X, Y] == \lambda R : X \leftrightarrow Y; T : \mathbb{P} Y \bullet \{p : R \mid p.2 \notin T\}$$

2.0.0.21. Name $_ \sim$ – Relational inversion

2.0.0.22. Definition

function($_ \sim$)

$$_ \sim [X, Y] == \lambda R : X \leftrightarrow Y \bullet \{p : R \bullet p.2 \mapsto p.1\}$$

2.0.0.23. Name $_ \langle _ \rangle$ – Relational image

2.0.0.24. Definition

function $(_ \circ _)$

$_ \circ _ [X, Y] == \lambda R : X \leftrightarrow Y; S : \mathbb{P} X \bullet \{p : R \mid p.1 \in S \bullet p.2\}$

2.0.0.25. Name Upper Bound

2.0.0.26. Definition

$upperBound[X, Y] == \lambda R : X \leftrightarrow Y \bullet \lambda S : \mathbb{P} X \bullet \{y : Y \mid S \times \{y\} \subseteq R\}$

2.0.0.27. Description The upper bound $upperBound\ R\ S$ of a set S through a relation R is the set of all objects y to which R relates every member x of S .

2.0.0.28. Name

- \circ – Relational composition
- \circ – Functional composition

2.0.0.29. Definition

function 40 leftassoc $(_ \circ _)$

$_ \circ _ [X, Y, Z] == \lambda Q : X \leftrightarrow Y; R : Y \leftrightarrow Z \bullet \{q : Q; r : R \mid q.2 = r.1 \bullet q.1 \mapsto r.$

function 40 leftassoc ($_ \circ _$)

$$_ \circ _[X, Y, Z] == \lambda R : Y \leftrightarrow Z; Q : X \leftrightarrow Y \bullet \{q : Q; r : R \mid q.2 = r.1 \bullet q.1 \mapsto r\}$$

2.0.0.30. Name *disjoint* – Disjointness
partition – Partitions

2.0.0.31. Definition

relation(*disjoint* $_$)

$$disjoint_ [I, X] == \{S : I \leftrightarrow \mathbb{P} X \mid \forall p, q : S \mid p \neq q \bullet p.2 \cap q.2 = \emptyset\}$$

relation($_ partition _$)

$$_ partition _[I, X] == \{S : I \leftrightarrow \mathbb{P} X; T : \mathbb{P} X \mid disjoint\ S \wedge \bigcup (ran\ S) = T\}$$

2.0.0.32. Name \oplus – Overriding

2.0.0.33. Definition

function 50 leftassoc ($_ \oplus _$)

$$_ \oplus _[X, Y] == \lambda Q, R : X \leftrightarrow Y \bullet (dom\ R \triangleleft Q) \cup R$$

2.0.0.34. Orders

2.0.0.35. Name

transitive	– Transitive relations
antisymmetric	– Antisymmetric relations
reflexive	– Reflexive relations
irreflexive	– Irreflexive relations
preOrder	– Preorders
order	– Partial orders
reflexiveOrder	– Reflexive partial orders
irreflexiveOrder	– Irreflexive partial orders
totalOrder	– Total orders
reflexiveTotalOrder	– Reflexive total orders
irreflexiveTotalOrder	– Irreflexive total orders
reflexiveChain	– Reflexive chains
irreflexiveChain	– Irreflexive chains

2.0.0.36. Definition

$\text{generic}(\text{transitive } _)$

$\text{transitive } X == \{R : X \leftrightarrow X \mid R \circ R \subseteq R\}$

$\text{generic}(\text{antisymmetric } _)$

$\text{antisymmetric } X == \{R : X \leftrightarrow X \mid R \cap R^\sim \subseteq \text{id } X\}$

`generic(reflexive _)`

$reflexive\ X == \{R : X \leftrightarrow X \mid id\ X \subseteq R\}$

`generic(irreflexive _)`

$irreflexive\ X == \{R : X \leftrightarrow X \mid R \cap id\ X = \emptyset\}$

`generic(preOrder _)`

$preOrder\ X == transitive\ X \cap reflexive\ X$

`generic(order _)`

$order\ X == transitive\ X \cap antisymmetric\ X$

`generic(reflexiveOrder _)`

$reflexiveOrder\ X == reflexive\ X \cap order\ X$

`generic(irreflexiveOrder _)`

$irreflexiveOrder\ X == irreflexive\ X \cap order\ X$

$generic(totalOrder\ _)$

$totalOrder\ X == \{R : order\ X \mid R \cup R^{\sim} \cup id\ X = X \times X\}$

$generic(reflexiveTotalOrder\ _)$

$reflexiveTotalOrder\ X == reflexive\ X \cap totalOrder\ X$

$generic(irreflexiveTotalOrder\ _)$

$irreflexiveTotalOrder\ X == irreflexive\ X \cap totalOrder\ X$

$reflexiveChain[X] == \lambda R : X \leftrightarrow X \bullet$
 $\{S : \mathbb{P}\ X \mid S \triangleleft R \triangleright S \in reflexiveTotalOrder\ S\}$

$irreflexiveChain[X] == \lambda R : X \leftrightarrow X \bullet$
 $\{S : \mathbb{P}\ X \mid S \triangleleft R \triangleright S \in irreflexiveTotalOrder\ S\}$

2.0.0.37. Name

- $_{}^+$ – Transitive closure
- $_{}^*$ – Reflexive-transitive closure
- do – Functional closure

2.0.0.38. Definition

function($_{-}^{+}$)

$$_{-}^{+}[X] == \lambda R : X \leftrightarrow X \bullet \bigcap \{Q : X \leftrightarrow X \mid R \subseteq Q \wedge R \circ Q \subseteq Q\}$$

function($_{-}^{*}$)

$$_{-}^{*}[X] == \lambda R : X \leftrightarrow X \bullet id\ X \cup R^{+}$$

$$do[X] == \lambda R : X \leftrightarrow X \bullet \bigcap \{Q : X \leftrightarrow X \mid id(X \setminus dom\ R) \subseteq Q \wedge R \circ Q \subseteq Q\}$$

3. Functions

3.0.0.39. Name

- \leftrightarrow – Partial functions
- \rightarrow – Total functions
- \rightsquigarrow – Partial injections
- \rightharpoonup – Total injections
- \twoheadrightarrow – Partial surjections
- \twoheadrightarrow – Total surjections
- \twoheadrightarrow – Bijections

3.0.0.40. Definition

generic 5 rightassoc $(_ \twoheadrightarrow _)$

$$X \twoheadrightarrow Y == \{f : X \leftrightarrow Y \mid \forall p, q : f \mid p.1 = q.1 \bullet p = q\}$$

generic 5 rightassoc $(_ \rightarrow _)$

$$X \rightarrow Y == \{f : X \twoheadrightarrow Y \mid \text{dom } f = X\}$$

generic 5 rightassoc $(_ \twoheadmap _)$

$$X \twoheadmap Y == \{f : X \leftrightarrow Y \mid \forall p, q : f \bullet p.1 = q.1 \Leftrightarrow p.2 = q.2\}$$

generic 5 rightassoc $(_ \mapsto _)$

$$X \mapsto Y == \{f : X \twoheadmap Y \mid \text{dom } f = X\}$$

generic 5 rightassoc $(_ \twoheadmapsto _)$

$$X \twoheadmapsto Y == \{f : X \twoheadrightarrow Y \mid \text{ran } f = Y\}$$

generic 5 rightassoc $(_ \twoheadrightarrow _)$

$$X \twoheadrightarrow Y == \{f : X \rightarrow Y \mid \text{ran } f = Y\}$$

generic 5 rightassoc $(_ \mapsto _)$

$$X \mapsto Y == \{f : X \mapsto Y \mid \text{ran } f = Y\}$$



Go Back

Full Screen

Close

Quit

- 3.0.0.41. Name**
- \rightarrow – Finite functions
 - \rightarrowtail – Finite injections

3.0.0.42. Definition

generic 5 rightassoc ($_{-} \rightarrow -$)

$$X \rightarrow Y == \{f : X \rightarrowtail Y \mid finite\ f\}$$

generic 5 rightassoc ($_{-} \rightarrowtail -$)

$$X \rightarrowtail Y == \{f : X \rightarrow Y \mid finite\ f\}$$

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