

normalization

[/Reference manual/Z-related commands/In situ replacement commands](#)

The *normalization* command extracts the property that is implicit in a declaration, reducing the declaration to be in terms of the carrier set of the underlying type (unlike *predication*).

$$\begin{aligned} \dots; i : e; \dots \mid p &\implies \dots; i : \text{carrier}(\tau e); \dots \mid i \in e \wedge p \\ \dots; i == e; \dots \mid p &\implies \dots; i : \text{carrier}(\mathbb{P}(\tau e)); \dots \mid i = e \wedge p \\ \dots; e; \dots \mid p &\implies \dots; \text{carrier}(\tau e); \dots \mid e \wedge p \end{aligned}$$

In the first two rules, it must be the name of the declaration that is inspected. In the last rule, that of a schema declaration, it must be the schema inclusion declaration that is inspected, not the schema expression within it—an extra click of button 1 is needed. The extracted property is that schema expression used as a predicate.

The *normalization* command can also be applied to declarations in the hypothesis part of a goal, the new predicate being generated as the first antecedent.

The *normalization* command can also be applied to a quantified predicate with a schema text whose \mid part is not *true*: the predicate of that \mid part is logically combined with the predicate in the \bullet part.

$$\begin{aligned} \forall ds \mid p_1 \bullet p_2 &\implies \forall ds \bullet p_1 \Rightarrow p_2 \\ \exists ds \mid p_1 \bullet p_2 &\implies \exists ds \bullet p_1 \wedge p_2 \\ \exists_1 ds \mid p_1 \bullet p_2 &\implies \exists_1 ds \bullet p_1 \wedge p_2 \end{aligned}$$



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See also [*denormalization*](#).

1. Tactic example

“normalization” d_1 d_2

This example applies the *normalization* command to declarations d_1 and d_2 .

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